



Computer Engineering Department

# Value-based Theoretical Guarantees

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Lecture 18 - 1

#### **Bellman's Optimality Equation**

• Assume a stochastic reward function.

$$\Pr(S_{t+1} = s', R_{t+1} = r | S_t = s, A_t = a), \forall s, s' \in \mathcal{S}, r \in \mathcal{R}, a \in \mathcal{A}, s' \in \mathcal{S}, r \in \mathcal{R}, s' \in \mathcal{S}, s' \in \mathcal{S}, r \in \mathcal{R}, s' \in \mathcal{S}, s' \in \mathcal$$

which is abbreviated by p(s', r|s, a).

$$\begin{aligned} q_*(s,a) &= \max_{\pi} \mathbb{E}[G_t | S_t = s, A_t = a] \\ &= \max_{\pi} \mathbb{E}[R_{t+1} + \gamma G_{t+1} | S_t = s, A_t = a] \\ &= \mathbb{E}[R_{t+1} | S_t = s, A_t = a] + \gamma \max_{\pi} \mathbb{E}[G_{t+1} | S_t = s, A_t = a]. \end{aligned}$$

# Bellman's Optimality Equation (cont.)

$$\mathbb{E}[R_{t+1}|S_t = s, A_t = a] = \sum_r r \sum_{s'} p(s', r|s, a).$$

$$\begin{split} \mathbb{E}[G_{t+1}|S_t = s, A_t = a] &= \sum_{s',a'} p(s',a'|s,a) \mathbb{E}[G_{t+1}|S_{t+1} = s', A_{t+1} = a', S_t = s, A_t = a] \\ &= \sum_{s',a'} p(s'|s,a) p(a'|s',s,a) \mathbb{E}[G_{t+1}|S_{t+1} = s', A_{t+1} = a'] \\ &= \sum_{s',a'} p(s'|s,a) \pi(a'|s') q_{\pi}(s',a') \\ &= \sum_{s'} p(s'|s,a) \sum_{a'} \pi(a'|s') q_{\pi}(s',a'). \end{split}$$

# Bellman's Optimality Equation (cont.)

$$q_*(s,a) = \sum_r r \sum_{s'} p(s',r|s,a) + \gamma \max_{\pi} \sum_{s'} p(s'|s,a) \sum_{a'} \pi(a'|s') q_{\pi}(s',a').$$

$$q_*(s,a) = \sum_r r \sum_{s'} p(s',r|s,a) + \gamma \max_{\pi} \sum_{s'} p(s'|s,a) \max_{a'} q_{\pi}(s',a').$$

# Bellman's Optimality Equation (cont.)

$$q_*(s,a) = \sum_{r} r \sum_{s'} p(s',r|s,a) + \gamma \sum_{s'} p(s'|s,a) \max_{a'} q_*(s',a')$$
$$= \sum_{r,s'} p(s',r|s,a)(r+\gamma \max_{a'} q_*(s',a')).$$

## Questions

- Does there exist q \* functions satisfying the Bellman's Eq.?
- Is this function unique?
- Can value iteration find this function?

## **Fixed Point**

- For an operator *T*, we call *x* a fixed point if Tx = x.
- $q_*$  is a fixed point of the Bellman's Eq.
- Why?

### Fixed Point (cont.)

**Theorem 1** (Banach Fixed Point Theorem). Suppose that X is a nonempty complete metric space and  $T: X \to X$  is a contraction mapping on X. Then T has a unique fixed point.

**Definition 1** (Contraction Mapping). [1] Let (X, d) be a metric space. A mapping  $T: X \to X$  is called a *contraction mapping* on X if there is a positive real number  $\alpha < 1$  such that for any  $x, y \in X$ 

 $d(Tx, Ty) \le \alpha d(x, y).$ 

#### **Existence** Proof

- Pick an arbitrary point x<sub>0</sub>.
- Construct a sequence: x

$$x_k = Tx_{k-1}, k = 1, 2, \dots$$

- Let  $C = d(x_1, x_0)$ .
- Note that

 $d(x_{k+1}, x_k) \le \alpha d(x_k, x_{k-1}) \le \dots \le \alpha^k d(x_1, x_0) = \alpha^k C, \, \forall, k = 1, 2, \dots$  $d(x_m, x_n) \le \sum_{i=0}^{m-n-1} d(x_{n+i+1}, x_{n+i}).$  $d(x_m, x_n) \le \sum_{i=0}^{m-n-1} \alpha^{n+i} C = \alpha^n C \frac{1 - \alpha^{m-n}}{1 - \alpha} \le \alpha^n \frac{C}{1 - \alpha}.$ 

#### Existence Proof

- Thus for any  $\epsilon > 0$ , if  $N \ge \frac{\log \epsilon (1-\alpha) \log C}{\log \alpha}$  then  $d(x_m, x_n) \le \epsilon$ .
- Hence *x<sub>n</sub>* is a Cauchy sequence.
- Therefore, it converges to a point, let's call x.
- Now, we show that x is a fixed point of T.
- Note that:

$$d(Tx, x) \le d(Tx, x_k) + d(x_k, x) \le \alpha d(x, x_{k-1}) + d(x_k, x), \forall k = 1, 2, \dots$$
  
 $d(Tx, x) = 0,$ 

#### Uniqueness

- Proof by contradiction.
- Let x' be another such fixed point.
- Then,  $d(x,x') = d(Tx,Tx') \le \alpha d(x,x'),$
- Which is a contradiction.

## Application to the Bellman's Eq.

• Define the operator *T* as:

$$Tq(s,a) = \sum_{r,s'} p(r,s'|s,a)(r + \gamma \max_{a'} q(s',a')),$$

#### T in Bellman is contraction

**Lemma 1.** For a finite MDP, the mapping T in Eq. (10) is a contraction mapping.

*Proof.* We consider the complete metric space  $(\mathbb{R}^{|\mathcal{S}| \times |\mathcal{A}|}, d)$ , where  $d(q_1, q_2) = ||q_1 - q_2||_{\infty}$  for any  $p, q \in \mathbb{R}^{|\mathcal{S}| \times |\mathcal{A}|}$ . Then,

$$\begin{split} \|Tq_1 - Tq_2\|_{\infty} &= \max_{s,a} |Tq_1(s,a) - Tq_2(s,a)| \\ &= \gamma \max_{s,a} \sum_{r,s'} p(r,s'|s,a) |\max_{a'} q_1(s',a') - \max_{a'} q_2(s',a')| \\ &\leq \gamma \max_{s,a} \sum_{s'} p(s'|s,a) \max_{a'} |q_1(s',a') - q_2(s',a')| \\ &\leq \gamma \max_{s,a} \max_{s'} \max_{a'} |q_1(s',a') - q_2(s',a')| \\ &= \gamma \max_{s',a'} |q_1(s',a') - q_2(s',a')| \\ &= \gamma \|q_1 - q_2\|_{\infty}, \end{split}$$

# Why value iteration converges to the fixed point?

• Let's discuss!

## Policy Improvement Improves!

- If we set the new policy to maximize q(s, a) over a, the new policy leads to higher v(s) values for all states s.
- Let's discuss!

#### **Policy Iteration Converges**

**Theorem.** Policy iteration is guaranteed to converge and at convergence, the current policy and its value function are the optimal policy and the optimal value function!

Proof sketch:

- (1) Guarantee to converge: In every step the policy improves. This means that a given policy can be encountered at most once. This means that after we have iterated as many times as there are different policies, i.e., (number actions)<sup>(number states)</sup>, we must be done and hence have converged.
- (2) Optimal at convergence: by definition of convergence, at convergence  $\pi_{k+1}(s) = \pi_k(s)$  for all states s. This means  $\forall s \ V^{\pi_k}(s) = \max_a \sum_{s'} T(s, a, s') \left[ R(s, a, s') + \gamma V_i^{\pi_k}(s') \right]$

Hence  $V^{\pi_k}$  satisfies the Bellman equation, which means  $V^{\pi_k}$  is equal to the optimal value function V<sup>\*</sup>.